

Quantum Search Algorithm

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1. Introduction

For a search problem associated with a unsorted database, the remarkable Grover's quantum algorithm (1) provides a quadratic speedup over its classical counterpart. The search problem can be described as follows: for a given function f , there exists one unknown element in the set $\{0, 1, \dots, N-1\}$ that satisfies $f(x) = -1$, say $x = \tau$, whereas the other $N-1$ ones give $f(x) = 1$. How many times of evaluations of f are required to determine the element τ for $f(\tau) = -1$? Through a conventional algorithm, one needs $O(N)$ trials to achieve this aim. How about the utility of quantum algorithm for the search?

For the scenario in the quantum world, the search problem can be rephrased in the quantum mechanical language: for a given unitary operator I_τ , that is sometimes called the *oracle operator*, and a set of state vectors (orthonormal basis): $\mathbf{s} = \{|0\rangle, |1\rangle, \dots, |N-1\rangle\}$, $I_\tau |x\rangle = |x\rangle$ for all states in the set except $I_\tau |x\rangle = -|x\rangle$ for $x = \tau$. How many queries of I_τ are required to determine $|\tau\rangle$? By Grover's algorithm, one needs only $O(\sqrt{N})$ quantum mechanical steps to find the marked state $|\tau\rangle$ out. It has been shown that Grover's algorithm is optimal since it needs minimal oracle calls to perform a quantum search (2).

The quantum searching process will be briefly reviewed as follows. The first step of Grover's algorithm is to prepare a superposition state of all elements with uniform probability amplitude:

$$|s\rangle = \frac{1}{\sqrt{N}} \sum_{x=0}^{N-1} |x\rangle. \quad (1)$$

Then apply the Grover kernel $G = -I_\eta I_\tau$ to $|s\rangle$, where I_η is a unitary operator and contains no bias against the marked state. For large N , after about $m = \pi\sqrt{N}/4$ repetitions of G operations, the probability to observe $|\tau\rangle$ is close to one, i.e.,

$$G^m |s\rangle \sim |\tau\rangle. \quad (2)$$

Since every single G involves one query of I_τ , only $O(\sqrt{N})$ searching steps are required for a quantum search task.

In what follows, we will first investigate on the general $SU(2)$ formulation for the kernel of Grover's searching operator G . The discussions of quantum searching certainty, robustness, and the analog analogue version of the Grover's algorithm will be given afterwards.

2. Kernel of Grover's searching operator G

Suppose in a two-dimensional complex Hilbert space we have a marked state $|\tau\rangle$ to be searched by successively operating a Grover's kernel G on an arbitrary initial state $|s\rangle$. The Grover kernel is a product of two unitary operators I_τ and I_η , given by (3)

$$I_\tau = I + (e^{i\phi} - 1) |\tau\rangle \langle \tau|, \quad (3)$$

$$I_\eta = I + (e^{i\theta} - 1) U |\eta\rangle \langle \eta| U^{-1}, \quad (4)$$

where U is an arbitrary unitary operator, $|\eta\rangle \in \mathbf{s}$, and ϕ and θ are two phase angles. Note that in the original designation of Grover, the phase angles $\phi = \theta = \pi$ are chosen for phase flips of the state vectors. The Grover kernel can be expressed in a matrix form in the following set of basis vectors:

$$|I\rangle = |\tau\rangle, \quad |II\rangle = (U |\eta\rangle - U_{\tau\eta} |\tau\rangle) / l, \quad (5)$$

where $U_{\tau\eta} = \langle \tau | U | \eta \rangle$ and $l = (1 - |U_{\tau\eta}|^2)^{1/2}$. Let us suppose that $U_{\tau\eta} = \sin(\beta)e^{i\alpha}$, we then have

$$U |\eta\rangle = \sin(\beta)e^{i\alpha} |I\rangle + \cos(\beta) |II\rangle, \quad (6)$$

and the Grover kernel can now be written as

$$\begin{aligned} G &= -I_\eta I_\tau \\ &= - \begin{bmatrix} e^{i\phi}(1 + (e^{i\theta} - 1) \sin^2(\beta)) & (e^{i\theta} - 1) \sin(\beta) \cos(\beta) e^{i\alpha} \\ e^{i\phi}(e^{i\theta} - 1) \sin(\beta) \cos(\beta) e^{-i\alpha} & 1 + (e^{i\theta} - 1) \cos^2(\beta) \end{bmatrix}. \end{aligned} \quad (7)$$

3. Quantum searching with certainty

In the searching process, the Grover kernel is successively operated on the initial state $|s\rangle$. We suppose that after m iterations the final state $G^m |s\rangle$ will be orthogonal to the basis vector $|II\rangle$, which means that the probability for finding the marked state $|\tau\rangle$ will exactly be unity, i.e.,

$$\langle II | G^m | s \rangle = 0, \quad (8)$$

which implies that

$$|\langle \tau | G^m | s \rangle| = |\langle I | G^m | s \rangle| = 1. \quad (9)$$

For the matrix of Eq. (7), the eigenvalues of the Grover kernel G are

$$\lambda_{1,2} = -e^{i(\frac{\phi+\theta}{2} \pm w)}, \quad (10)$$

where the angle w is defined by

$$\cos(w) = \cos\left(\frac{\phi - \theta}{2}\right) - 2 \sin\left(\frac{\phi}{2}\right) \sin\left(\frac{\theta}{2}\right) \sin^2(\beta). \quad (11)$$

Furthermore, the normalized eigenvectors associated with these eigenvalues can be derived:

$$|g_1\rangle = \begin{bmatrix} e^{-i\frac{\phi}{2}} e^{i\alpha} \cos(x) \\ \sin(x) \end{bmatrix}, \quad |g_2\rangle = \begin{bmatrix} -\sin(x) \\ e^{i\frac{\phi}{2}} e^{-i\alpha} \cos(x) \end{bmatrix}. \quad (12)$$

Here the angle x is defined by

$$\sin(x) = \sin\left(\frac{\theta}{2}\right) \sin(2\beta) / \sqrt{l_m}, \quad (13)$$

where

$$l_m = 2 \sin(w)(\sin(w) + \sin(\frac{\phi - \theta}{2})) + 2 \cos(\frac{\phi}{2}) \sin(\frac{\theta}{2}) \sin^2(\beta). \tag{14}$$

Then the matrix G^m can be expressed in the eigenbasis of G by $G^m = \lambda_1^m |g_1\rangle \langle g_1| + \lambda_2^m |g_2\rangle \langle g_2|$, and we have

$$G^m = (-1)^m e^{im(\frac{\phi+\theta}{2})} \begin{bmatrix} e^{imw} \cos^2(x) + e^{-imw} \sin^2(x) & e^{-i\frac{\phi}{2}} e^{i\alpha} i \sin(mw) \sin(2x) \\ e^{i\frac{\phi}{2}} e^{-i\alpha} i \sin(mw) \sin(2x) & e^{imw} \sin^2(x) + e^{-imw} \cos^2(x) \end{bmatrix}. \tag{15}$$

We proceed to consider the initial state $|s\rangle$. Here we assume that the initial state $|s\rangle$ is more general than the uniform superposition state (1):

$$|s\rangle = \sin(\beta_0) |I\rangle + \cos(\beta_0) e^{iu} |II\rangle. \tag{16}$$

Hence from the condition of searching with certainty, $\langle II|G^m|s\rangle = 0$, we have:

$$-\sin(mw) \sin(\frac{\phi}{2} - \alpha - u) \sin(2x) \sin(\beta_0) + \cos(mw) \cos(\beta_0) = 0, \tag{17}$$

$$\sin(mw) \cos(\frac{\phi}{2} - \alpha - u) \sin(2x) \sin(\beta_0) - \sin(mw) \cos(2x) \cos(\beta_0) = 0. \tag{18}$$

After substituting the angle x into Eq. (18), the equation can be reduced to the following condition:

$$(\sin(\frac{\phi - \theta}{2}) + 2 \cos(\frac{\phi}{2}) \sin(\frac{\theta}{2}) \sin^2(\beta)) \cos(\beta_0) = \sin(\frac{\theta}{2}) \sin(2\beta) \cos(\frac{\phi}{2} - \alpha - u) \sin(\beta_0), \tag{19}$$

which is identical to the relation derived by Long et al. (4):

$$\tan(\frac{\phi}{2}) = \tan(\frac{\theta}{2}) (\frac{\cos(2\beta) + \sin(2\beta) \tan(\beta_0) \cos(\alpha + u)}{1 - \tan(\beta_0) \tan(\frac{\theta}{2}) \sin(2\beta) \sin(\alpha + u)}). \tag{20}$$

For Eq. (17), under the satisfaction of the matching condition (19), or (20), we can have a formula for evaluating the number of iterations m :

$$\cos(mw + \sin^{-1}(\sin(\beta_0) \sin(\frac{\phi}{2} - \alpha - u))) = 0. \tag{21}$$

Then one can compute the number m :

$$m = \lceil f \rceil, \tag{22}$$

where $\lceil \cdot \rceil$ denotes the smallest integer greater than the quantity in it, and the function f is given by

$$f = \frac{\frac{\pi}{2} - \sin^{-1}(\sin(\beta_0) \sin(\frac{\phi}{2} - \alpha - u))}{\cos^{-1}(\cos(\frac{\phi - \theta}{2}) - 2 \sin(\frac{\phi}{2}) \sin(\frac{\theta}{2}) \sin^2(\beta))}. \tag{23}$$

It can be shown also that if the matching condition is fulfilled, then after m searching iterations the final state will be

$$G^m |s\rangle = e^{i\delta} |\tau\rangle = e^{i[m(\pi + \frac{\phi + \theta}{2}) + \Omega]} |\tau\rangle, \tag{24}$$

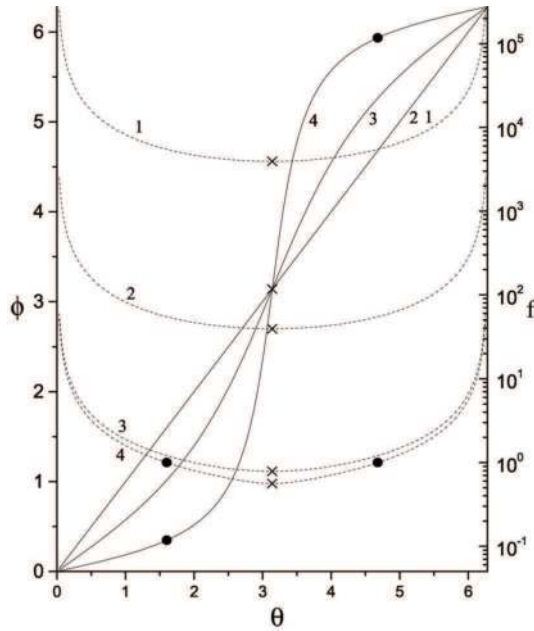


Fig. 1. Variations of $\phi(\theta)$ (solid) and $f(\theta)$ (broken), for $\alpha + u = 0, \beta_0 = 10^{-4}$, and $\beta = 10^{-4}$ (1), 10^{-2} (2), 0.5 (3) and 0.7 (4), respectively. The cross marks denote the special case of Høyer (6), while the black circles correspond to the optimal choices of ϕ_{op} and θ_{op} for $\alpha + u = 0, \beta_0 = 10^{-4}$ and $\beta = 0.7$. The solid straight line 1 corresponds the case $\phi = \theta$, while the solid curve 2 is only approximately close to the former.

where the angle Ω is defined by

$$\Omega = \tan^{-1}(\cot(\frac{\phi}{2} - \alpha - u)). \tag{25}$$

The matching condition (19), or (20), relates the angles $\phi, \theta, \beta, \beta_0$, and $\alpha + u$ for finding a marked state with certainty. If β, β_0 and $\alpha + u$ are designated, then $\phi = \phi(\theta)$ is deduced by the matching condition. As $\phi(\theta)$ is determined, we then can evaluate the value of $f = f(\phi(\theta), \theta)$ by (23) and consequently decide the number of iterations m by (22). It is worth to note that as $\alpha + u = 0$ and $\beta = \beta_0$, the matching condition recovers $\phi = \theta$ automatically since then Eq. (20) becomes an identity, and accordingly one has

$$f = \frac{\frac{\pi}{2} - \sin^{-1}(\sin(\frac{\phi}{2}) \sin(\beta))}{2 \sin^{-1}(\sin(\frac{\phi}{2}) \sin(\beta))}, \text{ for } \phi = \theta, \tag{26}$$

which is the case discussed in Ref. (5). We also note that the matching condition (19) will recover the relation by Høyer (6):

$$\tan(\frac{\phi}{2}) = \tan(\frac{\theta}{2}) \cos(2\beta), \text{ for } \cos(\phi/2 - \alpha - u) = 0. \tag{27}$$

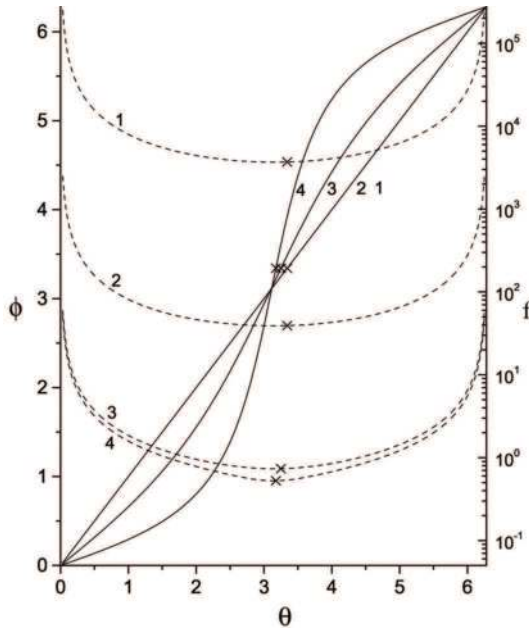


Fig. 2. Variations of $\phi(\theta)$ (solid) and $f(\theta)$ (broken), for $\alpha + u = 0.1$, $\beta_0 = 0.1$, and $\beta = 10^{-4}$ (1), 10^{-2} (2), 0.5 (3) and 0.7 (4), respectively. The cross marks denote the special case of Høyne (6). The solid curves 1 and 2 are almost identical, and both of them are only approximately close to the line $\phi = \theta$.

In Figs. 1 and 2 we have shown by the cross marks some particular examples of this special case.

Observing Figs. 1 and 2, one realizes that for every designation of β , β_0 and $\alpha + u$, the optimal choices for ϕ and θ is $\phi = \theta = \pi$, since then the corresponding f is minimum under the fact $df/d\theta = (\partial f/\partial\phi)(d\phi/d\theta) + \partial f/\partial\theta = 0$, for $\phi = \theta = \pi$. We thus denote the optimal value of m by

$$m_{op} = \lceil \min(f) \rceil = \left\lceil \frac{\frac{\pi}{2} - \sin^{-1}(\sin(\beta_0) \cos(\alpha + u))}{2\beta} \right\rceil. \tag{28}$$

With the choice of m_{op} , however, one needs to modify the phases θ and $\phi(\theta)$ to depart from π so that the matching condition is satisfied again. For example, if $\alpha + u = 0$, $\beta_0 = 10^{-4}$ and $\beta = 0.7$ are designated, then the minimum value of f will be $\min(f) = 0.56$. So we choose $m_{op} = 1$ and the modified phases are $\theta_{op} = (1 \pm 0.490)\pi$ and $\phi_{op} = (1 \pm 0.889)\pi$, respectively. This example has been shown by the marked black circles in Fig. 1. It is worth to note again that under the choice of m_{op} the modified ϕ and θ for the special case considered by Long (5) will be

$$\phi_{op} = \theta_{op} = \lceil \min(f) \rceil = 2 \sin^{-1} \left(\frac{\sin(\frac{\pi}{4m_{op}+2})}{\sin(\beta)} \right), \tag{29}$$

where

$$m_{op} = \left\lceil \frac{\frac{\pi}{2} - \beta}{2\beta} \right\rceil, \tag{30}$$

This is in fact a special case in which the phases ϕ_{op} and θ_{op} can be given by a closed-form formula.

4. Phase error tolerance in a quantum search algorithm

In Section 3 we have given a general review of quantum search algorithm with SU(2) formulation. To give a detailed discussion about phase error tolerance in a quantum search algorithm, in what follows we will focus on the original version of Grover's algorithm, which means that the unitary operator U shown in Eq. (4) corresponds to a specific transformation. Grover's quantum search algorithm (1) is achieved by applying Grover kernel G on an uniform superposition state (1), which is obtained by applying Walsh-Hadamard transformation on a initial state, in a specific operating step such that the probability amplitude of marked state is amplified to a desired one. Specifically, Grover's kernel is composed of phase rotations and Walsh-Hadamard transformations. The phase rotations include two kinds of operations: π -inversion of the marked state and π -inversion of the initial state. As shown in Section 3, the phases, π , can be replaced by two angles, ϕ and θ , under the phase matching criterion, which is the necessary condition for quantum searching with certainty. In other words, the relation between ϕ and θ will affect the degree of success of quantum search algorithm.

There have been several studies concern with the effect of imperfect phase rotations. Long *et al.* (7) have found that, a given expected degree of success P_{\max} , the tolerated angle difference between two phase rotations, δ , due to systematic errors in phase inversions is about $2/\sqrt{NP_{\max}}$, where N is the size of the database. Hoyer (6) has shown that after some number of iterations of Grover kernel, depending on N and unperturbed θ , it will give a solution with error probability $O(1/N)$ under a tolerated phase difference $\delta \sim O(1/\sqrt{N})$. The same result is also redrived by Biham *et al.* (8). A roughly close conclusion, $\delta O(1/N^{2/3})$, is presented by Pablo-Norman and Ruiz-Altaba (9).

The result of Long *et al.* (7) is based on the approximate Grover kernel and the assumptions of large N and small δ . However, one can find that the main inaccuracy comes from the approximate Grover kernel. Since all parameters in Grover kernel connect with each other exquisitely, any reduction to the structure of Grover's kernel would destroy this penetrative relation, so accumulative errors emerge from the iterations to a quantum search. In what follows, we will use the tool derived from SU(2) formulation discussed in Section 3 to get an improved criterion for tolerated error in phase rotation and the required number of qubits for preparing a database (10). In addition, a concise formula for evaluating minimum number of iterations to achieve a maximum probability will also be acquired.

Let us follow the derivation detailed in Section 3 for the analysis of error tolerance. If the operator U is Walsh-Hadamard transformation W , the orthonormal set becomes

$$|I\rangle = |\tau\rangle \text{ and } |\tau_{\perp}\rangle = (W|\eta\rangle - W_{\tau\eta}|\tau\rangle)/l, \quad (31)$$

where $W_{\tau\eta} = \langle\tau|W|\eta\rangle$ and $l = (1 - |W_{\tau\eta}|^2)^{1/2}$. Since $W_{\tau\eta} = \sin(\beta)$ we have

$$|s\rangle = W|\eta\rangle = \sin(\beta)|\tau\rangle + \cos(\beta)|\tau_{\perp}\rangle, \quad (32)$$

and the Grover kernel can now be written as

$$\begin{aligned} G &= -I_{\eta}I_{\tau} \\ &= - \begin{bmatrix} e^{i\phi}(1 + (e^{i\theta} - 1)\sin^2(\beta)) & (e^{i\theta} - 1)\sin(\beta)\cos(\beta) \\ e^{i\phi}(e^{i\theta} - 1)\sin(\beta)\cos(\beta) & 1 + (e^{i\theta} - 1)\cos^2(\beta) \end{bmatrix}. \end{aligned} \quad (33)$$

After m number of iterations, the operator G^m can be expressed as

$$G^m = (-1)^m e^{im(\frac{\phi+\theta}{2})} \begin{bmatrix} e^{imw} \cos^2(x) + e^{-imw} \sin^2(x) & e^{-i\frac{\phi}{2}} i \sin(mw) \sin(2x) \\ e^{i\frac{\phi}{2}} i \sin(mw) \sin(2x) & e^{imw} \sin^2(x) + e^{-imw} \cos^2(x) \end{bmatrix}. \quad (34)$$

Then the probability of finding a marked state is

$$\begin{aligned} P &= 1 - |\langle \tau | G^m | s \rangle|^2 \\ &= 1 - (\cos(mw) \cos(\beta) - \sin(mw) \sin(\frac{\phi}{2}) \sin(2x) \sin(\beta))^2 \\ &\quad - \sin^2(mw) (\cos(\frac{\phi}{2}) \sin(2x) \sin(\beta) - \cos(2x) \cos(\beta))^2. \end{aligned} \quad (35)$$

By the equation $\partial P / \partial (\cos(mw)) = 0$, the minimum number of iterations for obtaining the maximum probability, $P_{\max}(\cos(m_{\min}w))$, is evaluated:

$$m_{\min}(\beta, \phi, \theta) = \frac{\cos^{-1}(\sqrt{\frac{b-2a}{2b}})}{w}, \quad (36)$$

where

$$\begin{aligned} a &= \sin(2x) \cos(2\beta) + \cos(2x) \cos(\frac{\phi}{2}) \sin(2\beta), \\ b &= (2 + \sin^2(2x) + (3 \sin^2(2x) - 2) \cos(4\beta) - 2 \sin^2(2x) \cos(\phi) \sin^2(2\beta)) \\ &\quad + 2 \sin(4x) \cos(\frac{\phi}{2}) \sin(4\beta). \end{aligned}$$

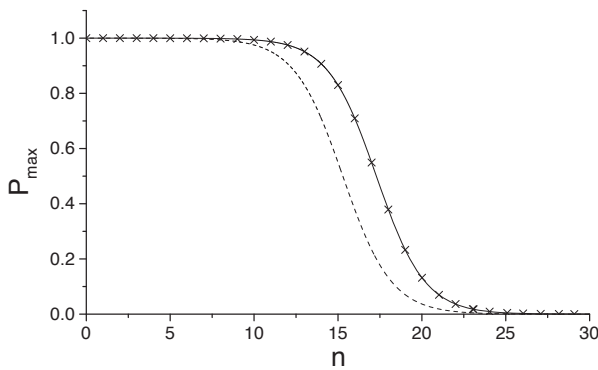


Fig. 3. Variations of exact vaule of $P_{\max}(n)$ (cross marks), $16\beta^2 \sin^2(\frac{\theta}{2}) / (\delta^2 + 16\beta^2 \sin^2(\frac{\theta}{2}))$ (solid), and $4\beta^2 / (\delta^2 + 4\beta^2)$ (dash) for $\theta = \pi, \delta = 0.01$ where $\beta = \sin^{-1}(2^{-n/2})$.

For a sure-success search problem, the phase condition, $\phi = \theta$, with $m_{\min} = (\pi/2 - \sin^{-1}(\sin(\phi/2) \sin(\beta))) / w$, is required. However, when effects of imperfect phase inversions are considered, the search is not certain, then the new condition to phase error, $\delta = \phi - \theta$, and the size of database should be derived in order to accomplish the search with a reduced maximum probability. Now, we suppose the database is large, i.e., if $\sin(\beta) \ll 1$, and a phase error δ is small, where $|\delta| \ll 1$, one has the following approximation:

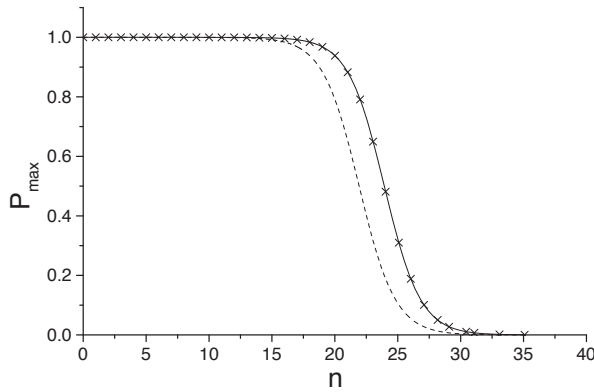


Fig. 4. Variations of exact value of $P_{\max}(n)$ (cross marks), $16\beta^2 \sin^2(\frac{\theta}{2}) / (\delta^2 + 16\beta^2 \sin^2(\frac{\theta}{2}))$ (solid), and $4\beta^2 / (\delta^2 + 4\beta^2)$ (dash) for $\theta = \pi, \delta = 0.001$ where $\beta = \sin^{-1}(2^{-n/2})$.

$$\begin{aligned} \cos(w) &= \cos(\frac{\delta}{2}) - 2 \sin(\frac{\theta}{2} + \frac{\delta}{2}) \sin(\frac{\theta}{2}) \sin^2(\beta) \\ &\approx 1 - (\frac{\delta^2}{8} + 2\beta^2 \sin^2(\frac{\theta}{2})), \\ \sin(w) &= (1 - \cos^2(w))^{1/2} \\ &\approx \frac{(\delta^2 + 16\beta^2 \sin^2(\frac{\theta}{2}))^{1/2}}{2}, \\ \sin(2x) &= \frac{4\beta \sin(\frac{\theta}{2})}{(\delta^2 + 16\beta^2 \sin^2(\frac{\theta}{2}))^{1/2}}. \end{aligned}$$

The probability P (Eq.35) then can be expressed approximately as

$$\begin{aligned} P &\approx 1 - \cos^2(mw) \cos^2(\beta) - \sin^2(mw) \cos^2(2x) \\ &= \sin^2(mw) \sin^2(2x), \end{aligned} \tag{37}$$

with a maximum value, by setting $\sin^2(mw) = 1$,

$$P_{\max} \approx \sin^2(2x) = \frac{16\beta^2 \sin^2(\frac{\theta}{2})}{\delta^2 + 16\beta^2 \sin^2(\frac{\theta}{2})}. \tag{38}$$

From Fig. 3 and Fig. 4, one can realize the function (38) depicted by solid line coincides with the exact value which is obtained by Eq. (35) and Eq. (36) as shown by cross marks. On the contrary, the result of Long *et al.*,

$$P_{\max} \approx \frac{4\beta^2 \sin^2(\frac{\theta}{2})}{\delta^2 + 4\beta^2 \sin^2(\frac{\theta}{2})}, \tag{39}$$

is an underestimation depicted by dash lines.

5. Family of sure-success quantum search algorithms

Even in the case of large N , where high success rate in finding the marked state is expected by using the standard Grover’s algorithm, inevitable noises including decoherence and gate inaccuracies can significantly affect the efficiency of the algorithm. To overcome such drawback , one might either apply the fault-tolerant computation (11) to reduce gate imperfections and decoherence, or limit the size of the quantum database to depress the effect of the uncertainty of the phase inversion operations. In another way, one can also, if possible, consider implementing a modified algorithm which is itself robust against phase imperfections and or decoherence. Recently, Hu (12) introduced an interesting family of algorithms for the quantum search. Although these algorithms are more complicated than the standard Grover’s algorithm, they can be proved to be robust against imperfect phase inversions, so the limitation of the size of database can be greatly relieved. In what follows the algorithms introduced by Hu (12) will be analyzed intently in detail, and the robustness of the family in resisting the effect of imperfect phase inversions will be shown later on.

Let us denote the phase inversion of marked state $I_\tau = 1 + (e^{i\phi} - 1) |\tau\rangle \langle\tau|$ as usual and the phase inversion of the initial state $I_s = 1 + (e^{i\theta} - 1) |s\rangle \langle s|$. Instead of applying G^n on the initial state $|s\rangle$, Hu (12) presented and utilized the operators $A_{2n} = (I_s^\dagger I_\tau^\dagger I_s I_\tau)^n$ and $A_{2n+1} = G A_{2n}$ to accomplish a quantum search with certainty, and named the former the even member and the latter the odd member of the family $\{A_n, n = 1, 2, \dots\}$ because they require even ($2n$) and odd ($2n + 1$) oracle calls in computation, respectively. The arrangement $I_s^\dagger I_\tau^\dagger I_s I_\tau$ is shown to have cancellation effect on phase errors in each iteration of the algorithm A_{2n} and A_{2n+1} and as a whole can ensure the robustness against imperfect phase inversion.

Consider a two-dimensional Hilbert space spanned by the marked state $|\tau\rangle$ and the state $|\tau_\perp\rangle$, which is orthogonal to $|\tau\rangle$. The initial state, as a uniform superposition of all states, then can be expressed by $|s\rangle = W|0\rangle = \sin(\beta) |\tau\rangle + \cos(\beta) |\tau_\perp\rangle$, where $\sin(\beta) \equiv \sqrt{M/N}$ and M is the number of the target states. The eigenvalues of the operator $I_s^\dagger I_\tau^\dagger I_s I_\tau$ are $\lambda_{1,2} = \cos(\omega) \pm i \sin(\omega)$ and the corresponding eigenvectors are computed (13; 14) as

$$\begin{aligned} |\lambda_1\rangle &= \cos(x) |\tau\rangle + i \sin(x) e^{i(\frac{\phi}{2} - \gamma)} |\tau_\perp\rangle, \\ |\lambda_2\rangle &= i \sin(x) e^{-i(\frac{\phi}{2} - \gamma)} |\tau\rangle + \cos(x) |\tau_\perp\rangle, \end{aligned} \tag{40}$$

where the rotation x and the related parameters are defined by

$$\tan(x) = \frac{2r \sin(\frac{\phi}{2}) \sin(\frac{\theta}{2}) \sin(2\beta)}{\sin(\omega) + \sin^2(\frac{\theta}{2}) \sin(\phi) \sin^2(2\beta)}, \tag{41}$$

$$\cos(\omega) = 1 - 2 \sin^2(\frac{\theta}{2}) \sin^2(\frac{\phi}{2}) \sin^2(2\beta), \tag{42}$$

$$r e^{i\gamma} = \cos(\frac{\theta}{2}) + i \sin(\frac{\theta}{2}) \cos(2\beta). \tag{43}$$

Then in n iterations of the operator of $I_s^\dagger I_\tau^\dagger I_s I_\tau$ we will have $A_{2n} = (I_s^\dagger I_\tau^\dagger I_s I_\tau)^n = \lambda_1^n |\lambda_1\rangle \langle\lambda_1| + \lambda_2^n |\lambda_2\rangle \langle\lambda_2|$, which can be expressed in the following matrix form:

$$A_{2n} = \begin{bmatrix} \cos(n\omega) + i \sin(n\omega) \cos(2x) & \sin(2x) \sin(n\omega) e^{-i(\frac{\phi}{2} - \gamma)} \\ -\sin(2x) \sin(n\omega) e^{i(\frac{\phi}{2} - \gamma)} & \cos(n\omega) - i \sin(n\omega) \cos(2x) \end{bmatrix}, \tag{44}$$

where $\sin(2x)$ and $\cos(2x)$ are given by

$$\sin(2x) = \left(\frac{1 - \sin^2(\frac{\theta}{2}) \sin^2(2\beta)}{1 - \sin^2(\frac{\theta}{2}) \sin^2(\frac{\phi}{2}) \sin^2(2\beta)} \right)^{1/2}, \quad (45)$$

$$\cos(2x) = \frac{\sin(\frac{\theta}{2}) \cos(\frac{\phi}{2}) \sin(2\beta)}{(1 - \sin^2(\frac{\theta}{2}) \sin^2(\frac{\phi}{2}) \sin^2(2\beta))^{1/2}}. \quad (46)$$

When the quantum search is carried out by using the even member A_{2n} , the component of the final state after n iterations of $(I_s^\dagger I_\tau^\dagger I_s I_\tau)$ in the basis state $|\tau_\perp\rangle$ is expressed by $\langle \tau_\perp | A_{2n} | s \rangle = RE_e + iIM_e$, and accordingly the exact success rate in finding the marked state $|\tau\rangle$ then is given by

$$p = 1 - |\langle \tau_\perp | A_{2n} | s \rangle|^2 = 1 - (RE_e^2 + IM_e^2), \quad (47)$$

where

$$RE_e = \cos(n\omega) \cos(\beta) - \sin(n\omega) \sin(2x) \cos(\frac{\phi}{2} - \gamma) \sin(\beta), \quad (48)$$

$$IM_e = - \frac{\sin(n\omega) \sin(\beta)}{(1 - \sin^2(\frac{\theta}{2}) \sin^2(\frac{\phi}{2}) \sin^2(2\beta))^{1/2}} \sin(\frac{\theta + \phi}{2}). \quad (49)$$

It is clear that when $IM_e = 0$, one obtains the n -independent phase matching condition, $\phi = -\theta$, for A_{2n} , and the success rate then becomes

$$p = 1 - RE_e^2 = 1 - \cos^2(n\omega - \alpha), \quad (50)$$

where $\alpha = \sin^{-1}(\sin(\beta) \cos(\phi/2 + \gamma))$. The 100% success rate for the search problem can be achieved as by setting $\cos(n\omega - \alpha) = 0$. For a search with certainty, since n is a positive integer, one therefore has to expect the iteration number given by

$$n_e(\theta, \beta) = \lceil f_e(\theta, \beta) \rceil, \quad (51)$$

and the function $f_e(\theta, \beta)$ is given by

$$f_e(\theta, \beta) = \frac{\frac{\pi}{2} + \alpha(\theta, \beta)}{\omega(\theta, \beta)}. \quad (52)$$

Given β , the function f_e has its minimal value as $\theta = \pi$ (and $\phi = -\pi$ thereby), as if minimal oracle calls are demanded in the computation, we should have the optimal phase θ_{op} associated with

$$f_e(\theta_{op}, \beta) = \lceil f_e(\pi, \beta) \rceil. \quad (53)$$

For example, if given $\beta = 1$, we have $\lceil f_e(\pi, 1) \rceil = 1$, and the optimal phase angle $\theta_{op} = \pi \pm 1.304$ follows in the algorithm using the even member A_{2n} . In usual operation, however, the quantum database is large, i.e., $\sin(\beta) \ll 1$, and the phase $\theta = \pi$ and $\phi = -\pi$ are fixed, then the required iterations are estimated by $n \sim \pi/8\beta$ and the maximal success rate will be approximately evaluated

$$p_{\max} \sim 1 - \beta^2, \text{ for } \theta = \pi,$$

which is the same result obtained as if the standard Grover algorithm is implemented. That is, as the phase $\theta = \pi$ is fixed, the present algorithm (A_{2n}) is equivalent to the standard algorithm (G^m) with even oracle calls required in the computation. Nevertheless, since in a real operation, imperfections in the phase inversions are inevitable. In what follows, one can

show that the present algorithm is robust against small phase imperfections in a quantum computation and provides a maximal success rate that is similar to the one given above. In the absence of decoherence and error correction, constant phase errors are considered to cause the phase ϕ and θ to be $\phi = \pi + \phi_e$ and $\theta = \pi + \theta_e$, where $|\phi_e| \ll 1$ and $|\theta_e| \ll 1$. By introducing the constant phase imperfections, one then has the following approximations, when $\beta \ll 1$,

$$\begin{aligned} \sin(2x) &\sim 1 - \frac{1}{2}\beta^2\phi_e^2, \cos(2x) \sim \beta\phi_e, \\ \cos\left(\frac{\phi}{2} - \gamma\right) &\sim -1 + \frac{1}{8}(\theta_e - \phi_e)^2, \sin\left(\frac{\phi}{2} - \gamma\right) \sim \frac{1}{2}(\theta_e - \phi_e), \\ \omega &\sim 4\beta\left(1 - \frac{1}{8}(\theta_e^2 + \phi_e^2) + \frac{4}{3}\beta^2\right). \end{aligned}$$

Then, since the errors are unknown in advance of the computation, the iteration number is also considered to be $n \sim \pi/8\beta$, and one thus has $\cos(n\omega) \sim \pi(\theta_e^2 + \phi_e^2)/16 - 2\pi\beta^2/3$ and $\sin(n\omega) \sim 1$. The approximation of RE_e and IM_e accordingly are evaluated by

$$\begin{aligned} RE_e &\sim \beta + \frac{\pi}{16}(\theta_e^2 + \phi_e^2) - \frac{2}{3}\pi\beta^2, \\ IM_e &\sim -\frac{1}{2}\beta(\theta_e + \phi_e). \end{aligned} \tag{54}$$

The maximal success rate now is approximately derived by

$$p_{\max} \sim 1 - \beta^2 - (\text{H.O.T.}), \tag{55}$$

where H.O.T. represents high order terms higher than second-degree in the small parameters β, θ_e and ϕ_e . Expression (55) clearly shows that the reduction of the probability due to the introduction of the phase errors in fact can almost be neglected. Then, through it, we can see that the present algorithm is robust against systematic phase imperfections.

We proceed to analyze the algorithm using the odd member A_{2n+1} by the same procedure as in analyzing the even member. In this case, one obtains

$$p_{\max} = 1 - (RE_o^2 + IM_o^2), \tag{56}$$

where

$$\begin{aligned} RE_o &= \cos(n\omega)\left[\cos\left(\frac{\theta - \phi}{2}\right) - 4\sin\left(\frac{\theta}{2}\right)\sin\left(\frac{\phi}{2}\right)\sin^2(\beta)\right]\cos(\beta) \\ &\quad + \sin(n\omega)\left\{\cos(2x)\left[\sin\left(\frac{\theta - \phi}{2}\right) - 4\sin\left(\frac{\theta}{2}\right)\sin\left(\frac{\phi}{2}\right)\sin^2(\beta)\right]\cos\beta\right. \\ &\quad \left. - \sin(2x)\left[\cos\left(\gamma + \frac{\theta}{2}\right) + 4\sin(\gamma)\sin\left(\frac{\theta}{2}\right)\cos^2(\beta)\right]\sin\beta\right\}, \\ IM_o &= \cos(\beta)\sin\left(\frac{\theta - \phi}{2}\right)(\cos(n\omega) - \sin(n\omega)\frac{\cos(2x)\sin\left(\frac{\phi}{2}\right)}{\cos\left(\frac{\phi}{2}\right)}). \end{aligned} \tag{57}$$

For $IM_o = 0$, one has the phase matching condition, $\phi = \theta$, for A_{2n+1} . The 100% success rate then can be ensured when the iteration steps at $n_o(\theta, \beta) = [f_o(\theta, \beta)]$, where

$$f_o(\theta, \beta) = \frac{\frac{\pi}{2} - \cos^{-1}\left(\frac{\cos(\beta)(1 - 4\sin^2(\frac{\theta}{2})\sin^2(\beta))\sqrt{1 - \sin^4(\frac{\theta}{2})\sin^2(2\beta)}}{\sqrt{1 - \sin^2(\frac{\theta}{2})\sin^2(2\beta)}}\right)}{\omega(\theta, \beta)}. \tag{58}$$

Note that in this case the inequality $1 - 4 \sin(\theta/2)^2 \geq 0$ should be demanded since then the meaningful requirement $f_o \geq 0$ can then be fulfilled. Given β , the function $f_o(\theta, \beta)$ also has its minimal value at $\theta = \pi$ (then $\phi = \pi$), as the optimal choice of the phase θ_{op} should be estimated by

$$f_o(\theta_{op}, \beta) = \lceil f_o(\theta, \beta) \rceil, \quad (59)$$

when minimal oracle calls are demanded in a search with certainty. For $\beta = 1$, the choice of the phase should be $\theta_{op} = \phi_{op} = \pi \pm 1.870$, for example. The standard Grover algorithm with odd oracle calls can be recovered when $\theta = \phi = \pi$ is fixed. In usual operations, when phase imperfections are introduced, i.e., as $\theta = \pi + \theta_o$ and $\phi = \pi + \phi_o$, where both θ_o and ϕ_o are small errors in the phases, they also produce almost negligible reductions in the success rate as given by an expression like Eq. (55).

6. Analog quantum search

Several researchers have proposed other ways to solve the quantum search problem, such as the *analog analogue* version of the Grover's algorithm (15–17) and the adiabatic evolution to quantum search (18–20). The former is to be considered here. It is proposed that the quantum search computation can be accomplished by controlled Hamiltonian time evolution of a system, obeying the Schrödinger equation

$$i \frac{d}{dt} |\Psi(t)\rangle = H |\Psi(t)\rangle, \quad (60)$$

where the constant $\hbar = 1$ is imposed for convenience. Farhi and Gutmann (15) presented the time-independent Hamiltonian $H_{fg} = E_{fg}(|w\rangle\langle w| + |s\rangle\langle s|)$, where $|w\rangle$ is the marked state and $|s\rangle$ denotes the initial state. Later, Fenner (16) proposed another Hamiltonian $H_f = E_f i(|w\rangle\langle s| - |s\rangle\langle w|)$. Recently, Bae and Kwon (17) further derived a generalized quantum search Hamiltonian

$$H_g = E_{fg}(|w\rangle\langle w| + |s\rangle\langle s|) + E_f(e^{i\phi}|w\rangle\langle s| + e^{-i\phi}|s\rangle\langle w|), \quad (61)$$

where ϕ is an additional phase to the Fenner Hamiltonian. Unlike the Grover algorithm, which operates on a state in discrete time, a quantum search Hamiltonian leads to the evolution of a state in continuous time, so the 100% probability for finding the marked state can be guaranteed in the absence of all kinds of imperfection occurring in a quantum operation. Both the Hamiltonian H_{fg} and H_f can help to find the marked state with 100% success. However, Bae and Kwon (17) addressed that the generalized Hamiltonian H_g can accomplish the search with certainty only when $\phi = n\pi$ is imposed, where n is arbitrary integer. In what follows, one can show that the generalized Hamiltonian H_g can be derived by an analytical method, which is distinct to the one implemented by Bae and Kwon (17), and the same method will lead to arbitrary chosen phase ϕ , depending on when the measurement on the system is undertaken and how large the system energy gap is provided. Since Hamiltonian-controlled system is considered, the energy-time relation will play an essential role in the problem. Therefore, the evaluation of the measuring time for the quantum search becomes crucially important. Here, the general Hamiltonian for the time-controlled quantum search system will be reviewed first. The exact time for measuring the marked state will be deduced. Finally, the role played by the phase ϕ in the quantum search will be discussed, and both the measuring time and the system energy gap as variations with ϕ will be given.

Suppose that a two-dimensional, complex Hilbert space is spanned by the orthonormal set $|w\rangle$, which is the marked state, and $|w_{\perp}\rangle$, which denotes the unmarked one. An initial state $|s\rangle = |\Psi(0)\rangle$, which corresponds to the quantum database discussed in the previous sections, is designed to evolve under a time-independent quantum search Hamiltonian given by (21)

$$H = E_1 |E_1\rangle \langle E_1| + E_2 |E_2\rangle \langle E_2|, \tag{62}$$

where E_1 and E_2 are two eigenenergies of the quantum system, $E_1 > E_2$, and $|E_1\rangle$ and $|E_2\rangle$ are the corresponding eigenstates satisfying the completeness condition $|E_1\rangle \langle E_1| + |E_2\rangle \langle E_2| = \mathbf{1}$. The eigenstates can be assumed by

$$\begin{aligned} |E_1\rangle &= e^{i\alpha} \cos(x) |w\rangle + \sin(x) |w_{\perp}\rangle, \\ |E_2\rangle &= -\sin(x) |w\rangle + e^{-i\alpha} \cos(x) |w_{\perp}\rangle. \end{aligned} \tag{63}$$

where x and α are two parameters to be determined later based on the required maximal probability for measuring the marked state. Then the Hamiltonian can be written in the matrix form

$$H = \begin{bmatrix} E_p + E_o \cos(2x) & E_o \sin(2x) e^{i\alpha} \\ E_o \sin(2x) e^{-i\alpha} & E_p - E_o \cos(2x) \end{bmatrix}. \tag{64}$$

where $E_p = (E_1 + E_2)/2$ is the mean of eigenenergies and $E_o = (E_1 - E_2)/2$ represents half of the system energy gap. The major advantage of using the controlled Hamiltonian time evolution is that the marked state can always be searched with certainty in the absence of quantum imperfections. The crucial key of the present problem in turn is to decide when to measure the marked state by the probability of unity. So in what follows the relation between all the unknowns appearing in the system will be deduced in detail and the exact measuring time for finding the marked state with certainty will be evaluated later on.

The time evolution of the initial state is given by $|\Psi(t)\rangle = e^{-iHt} |s\rangle$. Therefore, the probability of finding the marked state will be $P = \left| \langle w | e^{-iHt} |s\rangle \right|^2 = 1 - \left| \langle w_{\perp} | e^{-iHt} |s\rangle \right|^2$. Without loss of generality, let us consider the problem of searching one target from N unsorted items. The general form of the initial state considered in this study is given by

$$|s\rangle = e^{iu} \sin(\beta) |w\rangle + \cos(\beta) |w_{\perp}\rangle, \tag{65}$$

where $\sin(\beta) \equiv 1/\sqrt{N}$ and u denotes the relative phase between the two components in the initial state. Note that the relative phase u may arise from a phase decoherence or an intended design during the preparation of the initial state. Now, because of $e^{-iHt} = e^{-iE_1t} |E_1\rangle \langle E_1| + e^{-iE_2t} |E_2\rangle \langle E_2|$, one can deduce

$$\begin{aligned} \langle w_{\perp} | e^{-iHt} |s\rangle &= e^{-iE_p t} ((\cos(\beta) \cos(E_o t) - \sin(\alpha - u) \sin(2x) \sin(\beta) \sin(E_o t)) \\ &\quad + i(\cos(2x) \cos(\beta) - \cos(\alpha - u) \sin(2x) \sin(\beta)) \sin(E_o t)). \end{aligned} \tag{66}$$

To accomplish the quantum search with maximal probability, the time-independent term $(\cos(2x) \cos(\beta) - \cos(\alpha - u) \sin(2x) \sin(\beta))$ must be vanished and thus the unknown x can be determined by

$$\cos(2x) = \frac{\sin(\beta) \cos(\alpha - u)}{\cos(\gamma)}, \text{ or } \sin(2x) = \frac{\cos(\beta)}{\cos(\gamma)}, \tag{67}$$

where γ is defined by $\sin(\gamma) = \sin(\beta) \sin(\alpha - u)$. The probability for finding the marked state then becomes

$$\begin{aligned}
 P &= 1 - \left| \langle w_{\perp} | e^{-iHt} | s \rangle \right|^2 \\
 &= 1 - \frac{\cos^2(\beta)}{\cos^2(\gamma)} \cos^2(E_0 t + \gamma).
 \end{aligned}
 \tag{68}$$

If the size of database N is large, then $\gamma \ll 1$ and the marked state $|w\rangle$ will be measured at $t = \pi/(2E_0)$ by a probability $p = 1 - \tan^2 \gamma \sim 1$, according to (68). Expression (68) also indicates that, by setting $\cos^2(E_0 t + \gamma) = 0$, one can measure the marked state with unit probability, no matter how large N is, at the time instants

$$t_j = \frac{(2j - 1)\pi/2 - \sin^{-1}(\sin(\beta) \sin(\alpha - u))}{E_0}, \quad j = 1, 2, \dots
 \tag{69}$$

In what follows, let us only focus on the first instant $t_1 = (\pi/2 - \sin^{-1}(\sin(\beta) \sin(\alpha - u)))/E_0$. It is clear that a larger E_0 , or equivalently a larger system energy gap, will lead to a shorter time for measuring the marked state with certainty. Meanwhile, as can be seen in (68), the probability for measuring the marked state varies with time as a periodic function whose frequency is the Bohr frequency E_0/π , so a larger E_0 will also result in a more difficult control on the measuring time. In other words, the measuring time should be controlled more precisely for a higher Bohr frequency in the state evolution since then a small error in the measuring time will cost a serious drop of the probability. However, the energy gap E_0 depends on the size of database N , as will be mentioned later.

With the relations $\sin(2x)$ and $\cos(2x)$, the present Hamiltonian now can be written by

$$H = \begin{bmatrix} E_p + E_0 \frac{\sin(\beta) \cos(\alpha - u)}{\cos(\gamma)} & E_0 \frac{\cos(\beta)}{\cos(\gamma)} e^{i\alpha} \\ E_0 \frac{\cos(\beta)}{\cos(\gamma)} e^{-i\alpha} & E_p - E_0 \frac{\sin(\beta) \cos(\alpha - u)}{\cos(\gamma)} \end{bmatrix},
 \tag{70}$$

which is represented in terms of the energies E_p and E_0 and the phase α . Alternatively, if we let

$$\begin{aligned}
 E_{fg} &= \frac{(E_p - E_0 \frac{\sin(\beta) \cos(\alpha - u)}{\cos(\gamma)})}{\cos^2(\beta)}, \\
 E_f e^{i(\phi - u)} &= \frac{E_0}{\cos(\gamma)} e^{i(\alpha - u)} - E_{fg} \sin(\beta),
 \end{aligned}
 \tag{71}$$

or inversely,

$$\begin{aligned}
 E_p &= E_{fg} + E_f \cos(\phi - u) \sin(\beta), \\
 E_0 &= ((E_f \cos(\phi - u) + E_{fg} \sin(\beta))^2 + E_f^2 \sin^2(\phi - u) \cos^2(\beta))^{\frac{1}{2}},
 \end{aligned}
 \tag{72}$$

then the Hamiltonian can also be expressed by

$$H = \begin{bmatrix} E_{fg}(1 + \sin^2(\beta)) + 2E_f \cos(\phi - u) \sin(\beta) & e^{iu}(E_f e^{i(\phi - u)} + E_{fg} \sin(\beta)) \cos(\beta) \\ e^{-iu}(E_f e^{-i(\phi - u)} + E_{fg} \sin(\beta)) \cos(\beta) & E_{fg} \cos^2(\beta) \end{bmatrix},
 \tag{73}$$

which in turn is represented in terms of the energies E_{fg} and E_f and the phase ϕ . The Hamiltonian shown in (73) in fact can be expressed as $H_g = E_{fg}(|w\rangle\langle w| + |s\rangle\langle s|) + E_f(e^{i\phi}|w\rangle\langle s| + e^{-i\phi}|s\rangle\langle w|)$, which is exactly of the same form as the Bae and Kwon Hamiltonian H_g shown in (61). However, Bae and Kwon (17) only consider the case $u = 0$. In both the presentations (70) and (73) of the Hamiltonian H , the corresponding measuring time for finding the marked state $|w\rangle$ with certainty is at

$$t_1 = \frac{\frac{\pi}{2} - \sin^{-1}(\sin(\beta) \sin(\alpha - u))}{E_o} = \frac{\frac{\pi}{2} - \sin^{-1}\left(\frac{E_f \sin(\beta) \sin(\phi - u)}{((E_f \cos(\phi - u) + E_{fg} \sin(\beta))^2 + E_f^2 \sin^2(\phi - u))^{\frac{1}{2}}}\right)}{((E_f \cos(\phi - u) + E_{fg} \sin(\beta))^2 + E_f^2 \sin^2(\phi - u) \cos^2(\beta))^{\frac{1}{2}}}. \quad (74)$$

Equation (74) indicates that when the phase difference $\alpha - u$, or $\phi - u$, is imposed and the energy gap E_o or the energies E_f and E_{fg} are provided, the measurement at the end of a search should be undertaken at the instant t_1 . To discuss further, one first considers the case $u = 0$, i.e., the case where neither phase decoherence nor intended relative phase is introduced in the preparation of the initial state $|s\rangle$. If $\phi = n\pi$, or $\alpha = n\pi$, is imposed, then the present Hamiltonian reduces to that considered by Bae and Kwon (17) to serve for a search with certainty when the measurement is undertaken at $t_1 = \pi/(2E_o) = \pi/(2|(-1)^n E_f + E_{fg} \sin(\beta)|)$. If $E_f = 0$, or if $E_o = E_{fg} \sin(\beta)$ and $\alpha = 0$, is imposed, then the present Hamiltonian reduces to the Farhi and Gutmann Hamiltonian H_{fg} , which serves for a search with certainty at $t_1 = \pi/(2E_o) = \pi/(2E_{fg} \sin \beta)$. Further, when $E_{fg} = 0$ and $\phi = \pi/2$, or $E_p = 0$ and $\alpha = \pi/2$ is chosen, the present Hamiltonian will reduce to the Fenner Hamiltonian H_f associated with the measuring time $t_1 = (\pi - 2\beta)/(2E_o) = (\pi - 2\beta)/(2E_f \cos \beta)$. In general, the phase ϕ , or α , in fact can be imposed *arbitrary* for a search with certainty as the condition $u = 0$ is imposed.

However, if inevitable phase decoherence in the preparation of the initial state $|s\rangle$ is considered, then the phase u must be assumed to be arbitrary. Accordingly, the probability for finding the marked state will not be unity at all. For example, if following the treatment of Bae and Kwon (17) by setting $t_1 = \pi/(2E_o)$, then we only have a probability for finding the marked state given by

$$p = 1 - \frac{\cos^2(\beta) \sin^2(\beta) \sin^2(u)}{1 - \sin^2(\beta) \sin^2(u)}. \quad (75)$$

It is easy to show that the probability shown in (75) is always greater than or equal to the lower bound $p_{min} = 1 - \sin^2(\beta) = 1 - 1/N$. Of course, if the nonzero phase u is introduced by an intended design, not an inevitable phase decoherence, then a search with certainty can be accomplished for an arbitrary ϕ , or α , when associated with the measuring time shown in (74). For example, if $u = \pi/2$ is the phase designated, the ideal measuring time should be $t_1 = (\pi - 2\beta)/(2E_o)$, which is the same as the Fenner's t_1 . Again if the phase decoherence is introduced into the system and changes the phase from $\pi/2$ to an undesired u , then one eventually obtains a poor probability $p = 1 - (1 + \sin^2(u) - 2\sin(u)) \sin^2(\beta)$. Moreover if the phase error occurs randomly in a quantum database, then one is unable to be sure when to take a measurement, and the probability for finding the marked state even drops off seriously in some cases. For investigating the effect of the random uncontrollable parameter u on p at a fixed measuring time, one has to average over all possible values of $p(\beta, u)$ about all

arbitrary values of phase parameter u . Fig. 5 shows the variation of the mean probability \bar{p} with β for cases of Bae-Kwon, Farhi-Gutmann and Fenner at the specific measuring times, $t_{1,BK} = t_{1,FG} = \pi/(2E_0)$ and $t_{1,F} = (\pi - 2\beta)/(2E_0)$, those Hamiltonian suggest in such a case. The same character of their proposals is that \bar{p} is sensitive to a phase decoherence as the database is small. The mean success probabilities of Bae-Kwon and Farhi-Gutmann are the same and always greater than the one of Fenner. Then the Hamiltonians presented by Bae and Kwon, and Farhi and Gutmann are more robust against the phase decoherence than the one proposed by Fenner especially for low values of N .

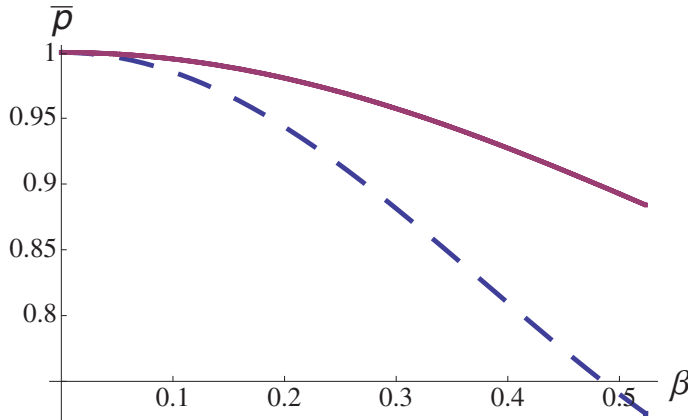


Fig. 5. The variation of $\bar{p}(\beta)$ for cases of Bae-Kwon(solid), Farhi-Gutmann(solid), and Fenner(broken) at the specific measuring times, $t_{1,BK} = t_{1,FG} = \pi/(2E_0)$ and $t_{1,F} = (\pi - 2\beta)/(2E_0)$.

A brief review on the comparison between E_{fg} and E_f , which has been discussed in Ref. (22), can be given now and the implication behind the analog quantum search first presented by Farhi and Gutmann (18) will be recalled. Suppose there is a $(N - 1)$ -fold degeneracy in a quantum system and its Hamiltonian is read as $H_0 = E |w\rangle \langle w|$, then our assignment is to find the unknown state $|w\rangle$. Since one does not yet know what $|w\rangle$ is, it is natural to add a well known Hamiltonian, $H_D = E |s\rangle \langle s|$, such that the initial state of the system $|s\rangle$ can be drove into $|w\rangle$. The total Hamiltonian therefore becomes $H = H_0 + H_D = E(|w\rangle \langle w| + |s\rangle \langle s|)$, which is just the Hamiltonian of Farhi and Gutmann (18) H_{fg} . It can be simplified under the large database limit,

$$H_{fg} \approx E(|w\rangle \langle w| + |s\rangle \langle s|) + E \sin(\beta)(|w\rangle \langle s| + |s\rangle \langle w|). \tag{76}$$

From it one can realize that the driving Hamiltonian induces transitions between $|w\rangle$ and $|s\rangle$ with a mixing amplitude $O(E \sin(\beta))$, which causes $|s\rangle$ to evolve to $|w\rangle$. By Eq. (72), thus it is rational to assume $E_f \sim E_{fg} \sin(\beta)$, and therefore the energy gap E_0 should be proportional to $\sin \beta$, or $1/\sqrt{N}$. The measuring time then is easily found to be $t_1 \propto \sqrt{N}$ from Eq. (74). However, if consider the case $E_f \gg E_{fg}$, like the extreme situation considered by Fenner (16), then one encounters with $E_0 \sim E_f \cos(\phi - u)$ and accordingly the measuring time t_1 is independent of the size of database N . Therefore, in an usual case the assumption $E_f \sim E_{fg} \sin(\beta)$ is reasonable.

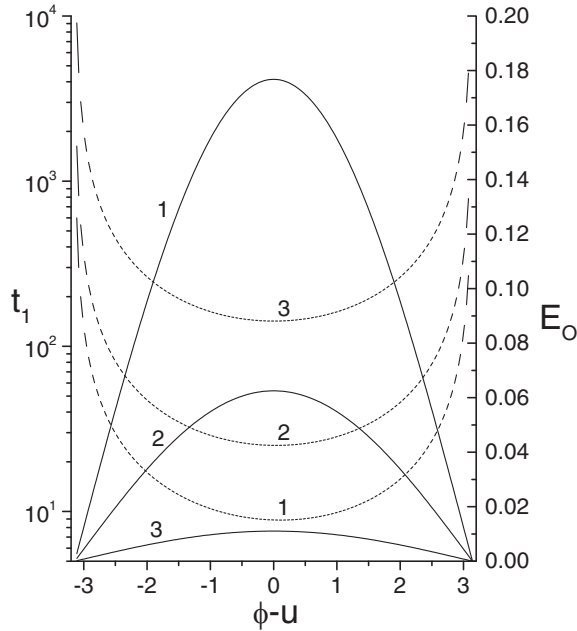


Fig. 6. Variations of $t_1(\phi - u)$ (broken) and $E_o(\phi - u)$ (solid), for $\beta = 0.085$ (1), $\beta = 0.031$ (2), and $\beta = 0.0055$ (3).

An interesting phenomenon occurs when the critical condition $E_f = E_{fg} \sin(\beta)$ is considered. Fig. 6 shows the variations of t_1 and E_o with the phase difference $\phi - u$ in such a case. It is observed that when $\phi - u = \pm\pi$ the energy gap E_o becomes zero and then the eigenstates of the quantum search system correspond to the common eigenvalue $E = E_1 = E_2$ and become degenerate. In such case, the Hamiltonian becomes proportional to the identity $\mathbf{1}(= |w\rangle \langle w| + |w_\perp\rangle \langle w_\perp|)$. Therefore, the initial state $|s\rangle$ does not evolve at all and the probability for finding the marked state $|w\rangle$ indeed is the initial one, viz., $p = \sin^2(\beta) = 1/N$, which can also be deduced using Eq. (68). In other words, the quantum search system is totally useless as long as $\phi - u = \pm\pi$ is imposed under the critical condition $E_f = E_{fg} \sin(\beta)$. When $\phi - u \neq \pm\pi$, both t_1 and E_o are finite, as can be seen from Fig. 6, and therefore the quantum search system becomes efficient again and is capable of finding the marked state with certainty, especially when the phase difference is imposed around $\phi - u = 0$. As a conclusion, for an efficient, useful quantum search system, the critical condition mentioned above should be avoided and in fact the reasonable condition $E_f \sim E_{fg} \sin(\beta)$ is recommended to be imposed.

7. Summary

A general SU(2) formulation is introduced to investigate the quantum search algorithm. In Sections 2 and 3, we show that the matching condition (19) for finding a marked state with certainty for arbitrary unitary transformations and an initial state can be derived from the general SU(2) formulation. Furthermore, one can also benefit from the same approach to evaluate the required number of iterations for the search such as Eqs. (22) and (23). With a given degree of maximum success, in Section 4 we derive a generalized and improved

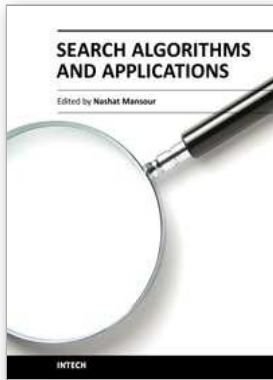
criterion for the tolerated error and the corresponding size of the quantum database under the inevitable gate imperfections. In Section 5, we consider a family of sure-success quantum algorithms and prove the matching conditions for both groups and give the corresponding formulae for evaluating the iterations or oracle calls required in the search. In addition, we show this kind of algorithms is robust against small phase imperfections in quantum gate operations. In the final section, we apply the same method as used for quantum search in the quantum circuit model to the analogue quantum search. A generalized Hamiltonian driving the evolution of quantum state in the analog search system is derived. Both the measuring time and the system energy gap suitable for a quantum search with or without certainty can be evaluated with the general results.

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Search algorithms aim to find solutions or objects with specified properties and constraints in a large solution search space or among a collection of objects. A solution can be a set of value assignments to variables that will satisfy the constraints or a sub-structure of a given discrete structure. In addition, there are search algorithms, mostly probabilistic, that are designed for the prospective quantum computer. This book demonstrates the wide applicability of search algorithms for the purpose of developing useful and practical solutions to problems that arise in a variety of problem domains. Although it is targeted to a wide group of readers: researchers, graduate students, and practitioners, it does not offer an exhaustive coverage of search algorithms and applications. The chapters are organized into three parts: Population-based and quantum search algorithms, Search algorithms for image and video processing, and Search algorithms for engineering applications.

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